

# Lecture 05: Chernoff Bound: An Introduction

# Introduction

- Let  $\mathbb{X}$  represent the Bern ( $p$ ) random variable
- Let  $\mathbb{X}^{(1)}, \dots, \mathbb{X}^{(n)}$  represent  $n$  independent and identical copies of the random variable  $\mathbb{X}$
- Let  $\mathbb{S}_n := \mathbb{X}^{(1)} + \dots + \mathbb{X}^{(n)}$  represent the sum of these  $n$  random variables
- Observe that  $\mathbb{S}_n$  is a random variable over the sample space  $\{0, 1, \dots, n\}$  with mean  $\mathbb{E}[\mathbb{S}_n] = np$
- For example, if  $\mathbb{X}$  represents a coin-toss, then  $\mathbb{S}_n$  is a random variable representing the number of observed Heads when  $n$  coin-tosses are performed
- How does the random variable  $\mathbb{S}_n$  concentrate around its mean? What is the probability of  $\mathbb{S}_n$  to be “far” from the expected value?

# Analysis using Markov Bound

- One can use Markov bound to deduce

$$\mathbb{P} [S_n \geq \lambda \cdot (np)] \leq \frac{1}{\lambda}.$$

- Can we do better?

- By Chebyshev's Inequality, we have

$$\mathbb{P} [|S_n - np| \geq t] \leq \frac{\text{Var}[S_n]}{t^2}.$$

- What is  $\text{Var} [S_n]$ ?

$$\begin{aligned}\text{Var} [S_n] &= \mathbb{E} [S_n^2] - (\mathbb{E} [S_n])^2 \\ &= \mathbb{E} \left[ \left( \sum_{i=1}^n X^{(i)} \right)^2 \right] - (np)^2 \\ &= \mathbb{E} \left[ \sum_{i=1}^n X^{(i)2} + \sum_{i \neq j} X^{(i)} X^{(j)} \right] - n^2 p^2 \\ &= n \cdot \mathbb{E} \left[ \left( X^{(1)} \right)^2 \right] + n(n-1) \cdot \mathbb{E} [X^{(1)} X^{(2)}] - n^2 p^2 \\ &= n \cdot p + n(n-1) \cdot p^2 - n^2 p^2 = n \cdot p(1-p).\end{aligned}$$

## Analysis using Chebyshev's Inequality III

- Think: The probability of  $S_n$  being  $\Theta\left(\sqrt{np(1-p)}\right)$  far from the mean is at most a constant.
- Think: Can we use higher moments to get better bounds?
- Think: Let  $(X_1, \dots, X_n)$  be a joint distribution and  $S_n = \sum_{i=1}^n X_i$ . Suppose the marginals  $X_i = \text{Bern}(p)$  and the random variables  $X_i$  and  $X_j$  are *pair-wise independent* when  $j \neq i$ . Can we still apply this estimation technique?

# A Large Deviation Bound

Observe that

$$\mathbb{P}[S_n \geq k] = \sum_{i=k}^n \binom{n}{i} \cdot p^i (1-p)^{n-i}.$$

Claim

$$\binom{n}{k} \cdot p^k (1-p)^{n-k} \leq \mathbb{P}[S_n \geq k] \leq \binom{n}{k} \cdot p^k.$$

- Think: How to prove this claim?
- Think: For what values of  $p$  and  $k$  is the upper bound meaningful? Hint: Use Stirling's formula.
- Think: When  $p = 1/2$ , for what values of  $k$  is the upper bound  $< 1$ ?

# Using Stirling's Approximation I

- Our objective is to study the expression

$$\mathbb{P}[S_n \geq k] = \sum_{i=k}^n \binom{n}{i} \cdot p^i (1-p)^{n-i}.$$

- It is obvious that this expression is at least the term  $\binom{n}{k} \cdot p^k (1-p)^{n-k}$ .
- In homework, we lower-bounded this term by

$$\frac{1}{\sqrt{8np'(1-p')}} \exp\left(-nD_{\text{KL}}(p', p)\right),$$

where  $p' = k/n$  and  $D_{\text{KL}}(a, b) = a \ln\left(\frac{a}{b}\right) + (1-a) \ln\left(\frac{1-a}{1-b}\right)$  represents the Kullback–Leibler divergence



## Using Stirling's Approximation II

- Therefore, to obtain a tight upper bound of the original expression, we should aim to obtain an upper bound that is in terms of  $\exp\left(-nD_{\text{KL}}(p', p)\right)$
- Towards this objective, we prove the following upper bound on the  $j$ -th term of the summation

### Claim

For  $j \geq 0$ , we have

$$\binom{n}{k+j} \cdot p^{k+j}(1-p)^{n-k-j} \leq \rho^j \cdot \binom{n}{k} \cdot p^k(1-p)^{n-k},$$

where  $\rho = \frac{1-p'}{p'} \cdot \frac{p}{1-p}$  and  $p' = k/n$ .

Think: Why is  $\rho < 1$ , if  $p' = k/n > p$ ?

# Using Stirling's Approximation III

- Therefore, we have

## Theorem

$$\begin{aligned}\mathbb{P}[S_n \geq k] &< \frac{1}{1-\rho} \cdot \binom{n}{k} \cdot p^k (1-p)^{n-k} \\ &\leq \frac{1}{1-\rho} \cdot \frac{1}{\sqrt{2\pi np'(1-p')}} \exp\left(-nD_{\text{KL}}(p', p)\right) \\ &= \frac{p'(1-p)}{(p'-p)} \cdot \frac{1}{\sqrt{2\pi np'(1-p')}} \exp\left(-nD_{\text{KL}}(p', p)\right),\end{aligned}$$

where  $\rho = \frac{1-p'}{p'} \cdot \frac{p}{1-p}$  and  $p' = k/n$ .

The first inequality follows from the claim above. The second inequality is an estimation using the Stirling's approximation, which we proved in homework.

# Using Stirling's Approximation IV

- Observe that if  $p'$  is a constant  $> p$ , then
  - ① The lower and the upper bounds are within a constant factor of each other!
  - ② The probability is exponentially decreasing in  $n$ .
- The conclusions are summarized in the next result

# Using Stirling's Approximation V

## Lemma (Conclusions)

Let  $S_n = \mathbb{X}^{(1)} + \dots + \mathbb{X}^{(n)}$ , where  $\mathbb{X} = \text{Bern}(p)$ .

1

$$\mathbb{P}[S_n \geq k] \leq \binom{n}{k} p^k.$$

2

$$1 \leq \frac{\mathbb{P}[S_n \geq k]}{\binom{n}{k} p^k (1-p)^{n-k}} \leq \frac{1}{1-\rho},$$

where  $\rho = \frac{1-p'}{p'} \cdot \frac{p}{1-p}$  and  $p' = k/n$ .

3

$$\binom{n}{k} p^k (1-p)^{n-k} = \frac{\Theta(1)}{\sqrt{np'(1-p')}} \cdot \exp\left(-n \cdot D_{\text{KL}}(p', p)\right),$$

where  $D_{\text{KL}}(a, b) = a \ln\left(\frac{a}{b}\right) + (1-a) \ln\left(\frac{1-a}{1-b}\right)$  and  $p' = k/n$ .